

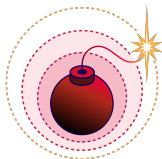
Blast-RADIUS

Breaking Enterprise Network Authentication

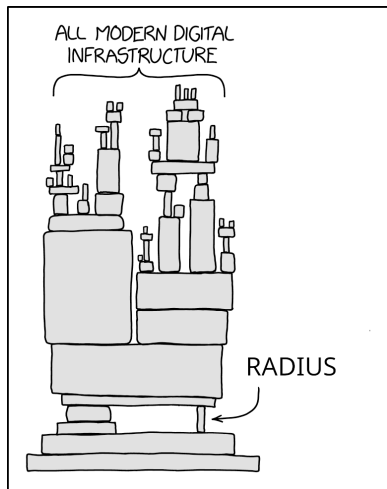
Sharon Goldberg¹, **Miro Haller**², **Nadia Heninger**², Mike Milano³, Dan Shumow⁴,
Marc Stevens⁵, **Adam Suhl**²

¹Cloudflare, ²**UC San Diego**, ³BastionZero, ⁴Microsoft Research, ⁵Centrum Wiskunde & Informatica

CNS 2025; May 9, 2025



What is RADIUS? Where is it used?

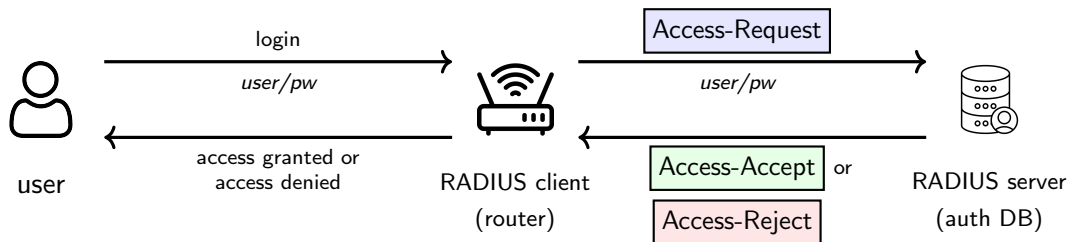


modified XKCD from [5]

- *RADIUS*: standard protocol for enterprise network authentication.
- RADIUS is *everywhere*:
RADIUS is [...] supported by essentially every switch, router, access point, and VPN concentrator product sold in the past twenty-five years.
(Alan DeKok [3])
- Used for backbone routers, non-cable ISP, IoT devices, identity providers (Okta, Duo), 802.1X, enterprise WiFi, eduroam...

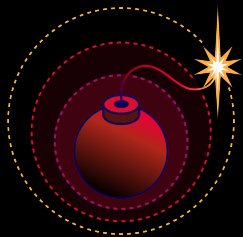
Blast-RADIUS on a Single Slide

How does RADIUS work?



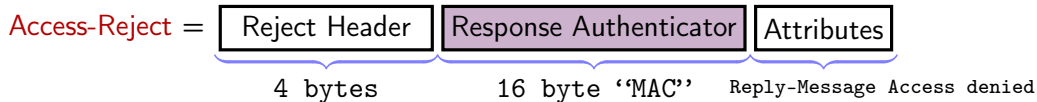
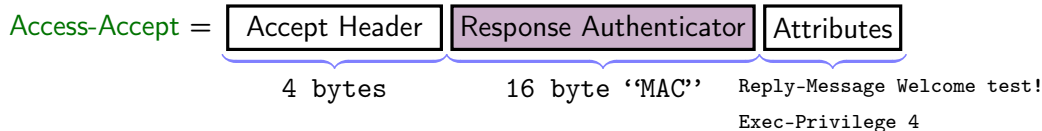
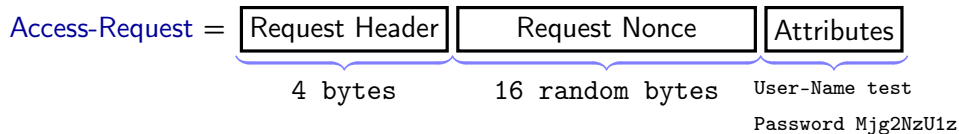
- RADIUS uses a custom MD5-based MAC and sends most traffic over UDP.
- *Our protocol vulnerability:* MITM can change Access-Reject to Access-Accept.
- *Impact:* authenticate as any user; accelerate RADIUS/UDP deprecation.
- *Mitigation:* responsible disclosure with over 90 vendors (incl. Cisco, Microsoft, ...).

icons from [4]



THE RADIUS PROTOCOL

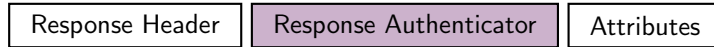
RADIUS Packet Formats



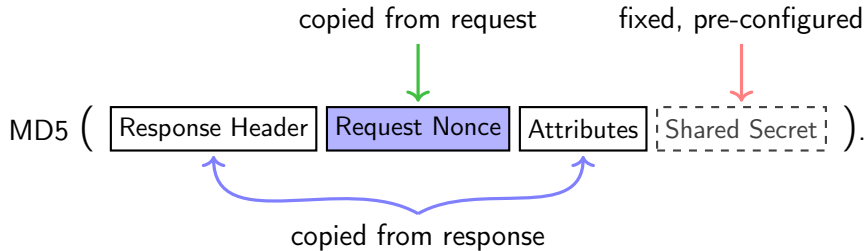
Response Authenticator

Goal: Prevent forgery of packets (e.g., by MITM attacker).

The Response Authenticator from packet



is computed as



90s Cryptography In RADIUS

RADIUS must be broken.



Sharon Goldberg

Let's do it!



Nadia Heninger

*As of the writing of this specification, RADIUS/UDP is still widely used, even though it depends on MD5 and "ad hoc" constructions for security. While MD5 has been broken, **it is a testament to the design of RADIUS that there have been (as yet) no attacks on RADIUS Authenticator signatures** which are stronger than brute-force.*

("Deprecating Insecure Practices in RADIUS" IETF draft, 2023)

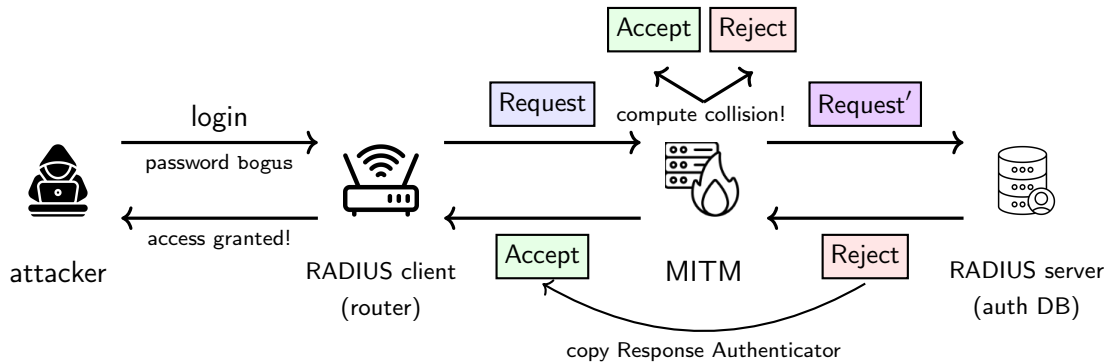


THE BLAST-RADIUS ATTACK

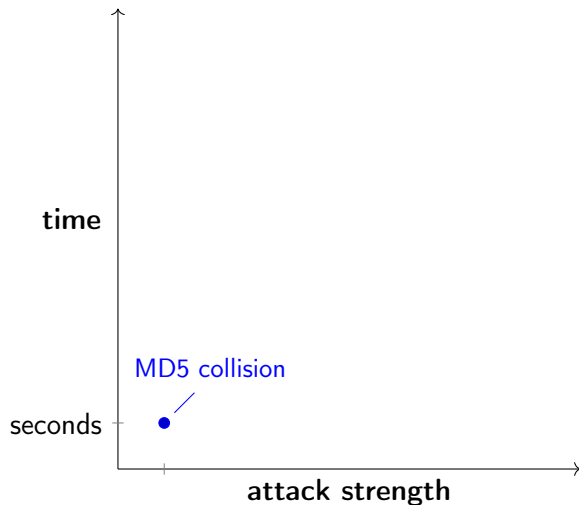
Blast-RADIUS: Attack Overview

Goal: Forge Access-Accept without knowing shared secret.

Blast-RADIUS attack: Create MD5 collision s.t. Access-Accept and Access-Reject produce same Response Authenticator: $\text{MD5}(\text{Access-Accept}) = \text{MD5}(\text{Access-Reject})$.



MD5 Collisions



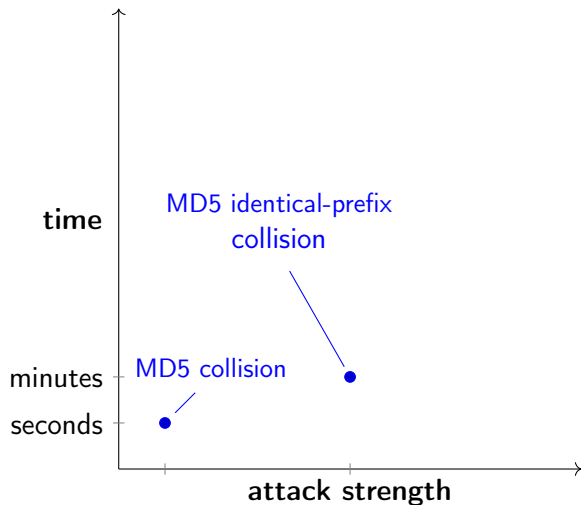
2004: MD5 collision [12]

Produce unstructured strings G_1 , G_2 such that

$$\text{MD5}(G_1) = \text{MD5}(G_2).$$



MD5 Collisions



2004: *Identical-prefix collision* [12]

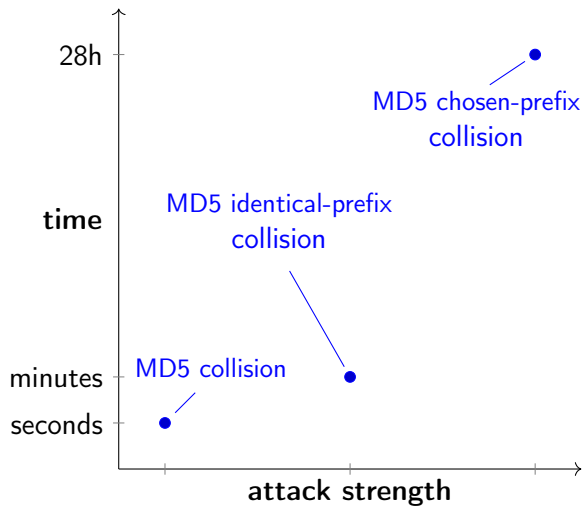
Given prefix P , produce G_1, G_2 such that

$$\text{MD5}(P||G_1) = \text{MD5}(P||G_2).$$



famous non-MD5 example of an identical-prefix collision [8]

MD5 Collisions



2007: MD5 chosen-prefix collision [9]

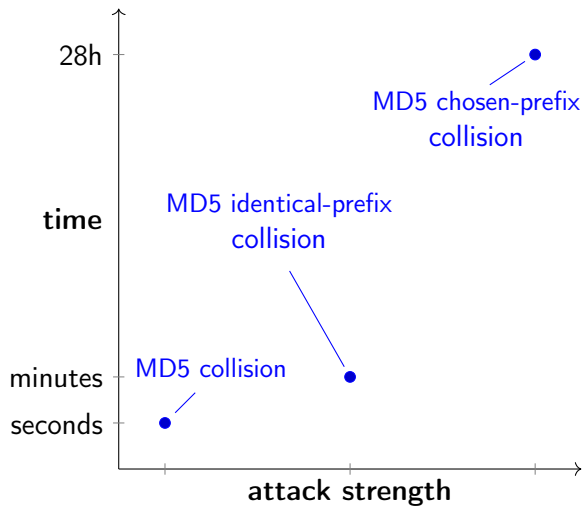
Given prefixes P_1 , P_2 , produce G_1 , G_2 such that

$$\text{MD5}(P_1 || G_1) = \text{MD5}(P_2 || G_2).$$



215 PS3 for Rogue TLS CA cert [10]

MD5 Collisions



2007: MD5 chosen-prefix collision [9]

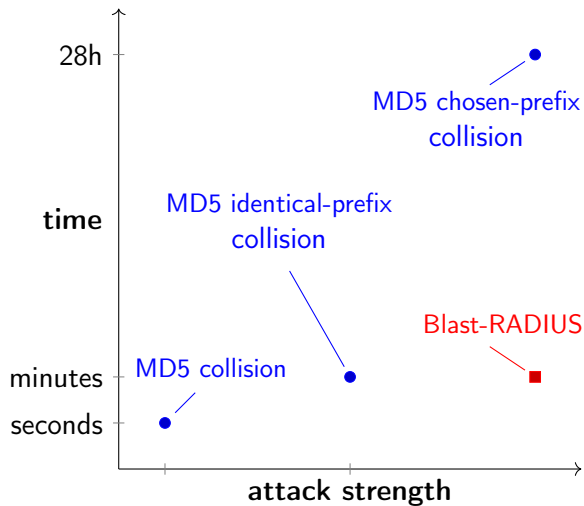
Given prefixes P_1 , P_2 , produce G_1 , G_2 such that

$$\text{MD5}(P_1 || G_1) = \text{MD5}(P_2 || G_2).$$

Due to Merkle-Damgård structure of MD5, can append identical suffix S :

$$\text{MD5}(P_1 || G_1 || S) = \text{MD5}(P_2 || G_2 || S).$$

MD5 Collisions



2024: *Blast-RADIUS*

A chosen-prefix collision:*

Given prefixes P_1 , P_2 , produce G_1 , G_2 such that (for any suffix S)

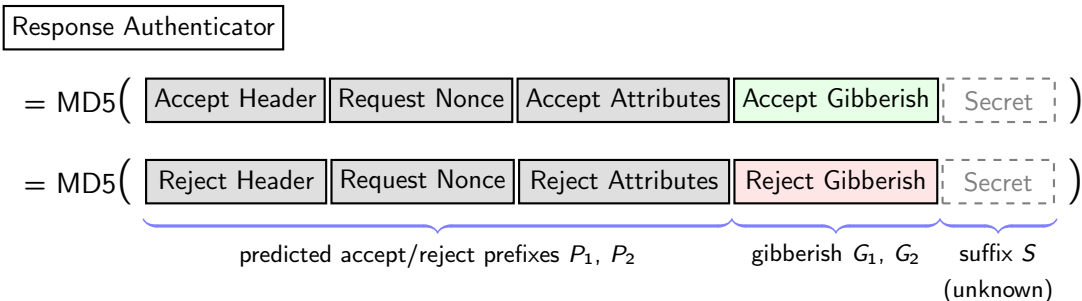
$$\text{MD5}(P_1 || G_1 || S) = \text{MD5}(P_2 || G_2 || S).$$

*But faster and with some additional conditions!

Blast-RADIUS: Turning Access-Reject into Access-Accept

Attack: MD5 collision to forge Access-Accept with same Response Authenticator as Access-Reject (without knowledge of the shared secret).

MD5 chosen-prefix collision $\text{MD5}(P_1 || G_1 || S) = \text{MD5}(P_2 || G_2 || S)$ applied to RADIUS:



That seems easy... too easy?



Challenge 1: Inject MD5 Reject Gibberish In Protocol

Problem: Server must include **Reject Gibberish** in Response Authenticator computation for Access-Reject.

$$\text{MD5}\left(\boxed{\text{Reject Header}} \boxed{\text{Request Nonce}} \boxed{\text{Reject Gibberish}} \boxed{\text{Shared Secret}} \right)$$

Solution: The Proxy-State attribute.

*This Attribute is available to be sent by a proxy server to another server when forwarding an Access-Request and **MUST be returned unmodified** in the Access-Accept, Access-Reject or Access-Challenge.*

(RFC 2058, emphasis added)



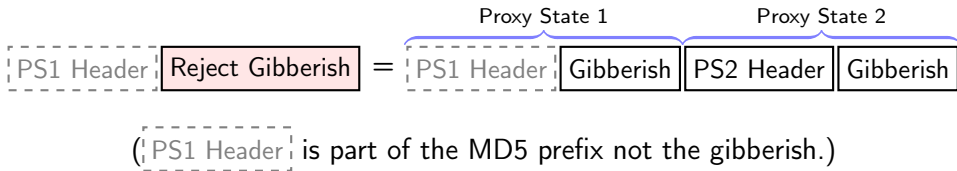
$$\begin{aligned} \text{Access-Request} &= \boxed{\text{Request Header}} \boxed{\text{Request Nonce}} \boxed{\text{Proxy-State}} \\ \text{Access-Reject} &= \boxed{\text{Reject Header}} \boxed{\text{Response Authenticator}} \boxed{\text{Proxy-State}} \end{aligned}$$

A blue arrow points from the Proxy-State box in the Access-Request line to the Proxy-State box in the Access-Reject line.

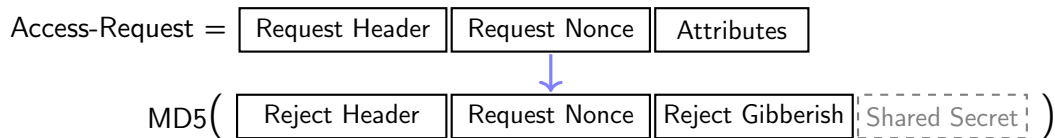
Challenge 2: Fitting Gibberish Into Protocol Attributes

Problem: Hiding Reject Gibberish in single Proxy-State attribute is too slow.

Solution: Spread longer gibberish across multiple Proxy-State attributes by modifying collision algorithm to embed Proxy-State header.



Challenge 3: Fast Collision Computation



Problem: Computing MD5 prefixes requires

Request Nonce

.

⇒ Must compute collision before RADIUS client times out.

Solution: Reduce collision time from days to $\leq 5\text{m}$ (on 47 servers) with algorithmic improvements and parallelization (next).

Challenge 3: Our Optimizations

- Increase precomputation.
- New GPU generator mode for birthday search.
- Parallelize forward phase across machines.
- Tune parameters to trade runtime for success rate.



Execution monitoring during collision optimization.

Blast-RADIUS: Example

As concrete example, putting everything together, we get the following collision.

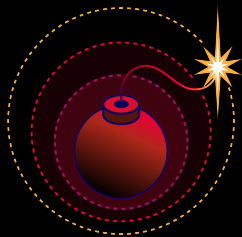
Response Authenticator

6034d0ff16e4...30

PoC example packets

blastradius.fail/example.py

$$\begin{aligned} &= \text{MD5}\left(\begin{array}{c} \text{Header} \quad \text{Request Nonce} \quad \text{Proxy State 1} \quad \text{Proxy State 2} \\ \boxed{02} \boxed{1d} \boxed{01c0} \boxed{726164617574...72} \boxed{21} \boxed{ec} \boxed{3d...86} \boxed{21} \boxed{c0} \boxed{f5...9e} \boxed{\text{Shared Secret}} \end{array} \right) \\ &\quad \text{Accept Prefix} \quad \text{Accept Gibberish} \quad \text{(unknown)} \\ \\ &= \text{MD5}\left(\begin{array}{c} \boxed{03} \boxed{1d} \boxed{01c0} \boxed{726164617574...72} \boxed{21} \boxed{ec} \boxed{96...86} \boxed{21} \boxed{c0} \boxed{f5...9e} \boxed{\text{Shared Secret}} \end{array} \right) \\ &\quad \text{Reject Prefix} \quad \text{Reject Gibberish} \quad \text{(unknown)} \end{aligned}$$



IMPACT

of the Blast-RADIUS attack

Successful PoCs*

Blast-RADIUS allows attacker to authenticate:

- **FreeRADIUS 3.2.3**: “most widely used RADIUS server in the world” [6]
- **Okta**: RADIUS in PAP mode for MFA.
- **Cisco ASA 5505 firewall** using RADIUS to authenticate users for access to serial console, VPN, Telnet, FTP, or HTTPS.
- **PAM**: RADIUS authentication for SSH, sudo.

⇒ Confirms no Message-Authenticator used, Proxy-State accepted in Access-Accept.

*With longer timeouts than used in practice.



PoC with Cisco ASA 5505 firewall tunneling UDP via TCP to our cluster.

Impact Summary

Affected authentication modes:

- PAP, CHAP, MS-CHAP are vulnerable.
- EAP not vulnerable (more later).

Affected deployments: Requires MITM network access

- RADIUS/UDP over Internet: vulnerable.
 - incl. many non-cable ISPs and major cloud providers.
- RADIUS/UDP over VLAN/IPsec: lateral movement.

Timing:

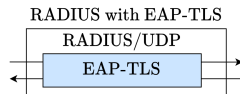
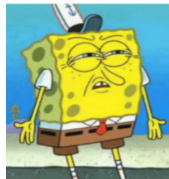
- RADIUS client timeouts $\leq 1\text{m}$, our PoCs take $\approx 5\text{m}$.
- Optimizations possible (parallelizes well, hardware implementation) but time consuming.

“Lots and lots and lots of people [are] sending [RADIUS/]UDP over the public Internet”

(Alan DeKok)

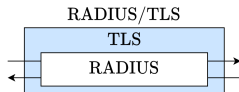
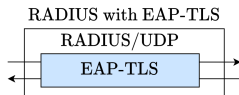
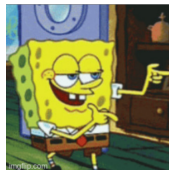
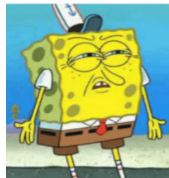
EAP: It's Complicated.

- TLS in EAP-TLS does not protect RADIUS packets.



EAP: It's Complicated.

- TLS in EAP-TLS does not protect RADIUS packets.
- Not to be confused with RADIUS/TLS, which properly nests RADIUS inside TLS.
- RFC 3579 requires that EAP-Message has Message-Authenticator attribute [1].
- Unclear client behavior for Access-Accept without EAP-Message.
- In eduroam and 802.1X, key is negotiated inside EAP session $\implies$ would require further attacks.





MITIGATING

the Blast-RADIUS attack

Mitigations

Massive disclosure with 90+ vendors.

Long-term:

- Encapsulate all RADIUS traffic in (D)TLS tunnel.
- Current IETF draft is being standardized [7].

Challenges: widely used, need backwards compatibility.

Short-term:

- Message-Authenticator attribute uses HMAC-MD5 not vulnerable to MD5 collisions.
- All requests and responses should include and verify Message-Authenticator.

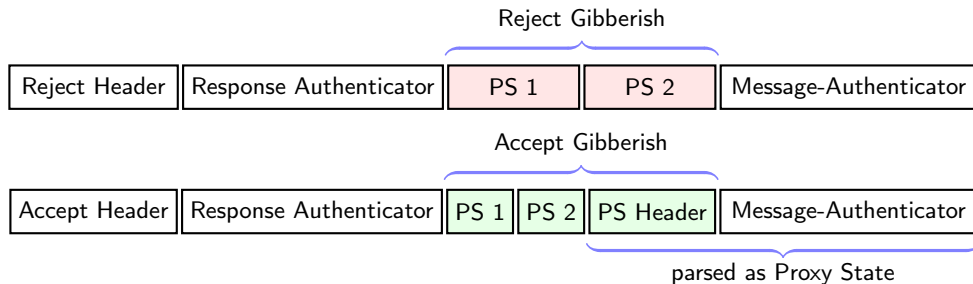


Some power plants use RADIUS [11].

Mitigation Troubles

Many equipment vendors have upgraded [2], but some challenges remain:

- Juniper vs. Cisco: incompatible Message-Authenticator placement.
- Correct behavior: put as first attribute for sending, mandate presence for receiving.
- Incorrect placement may be vulnerable to **Message-Authenticator hiding attack**:



Blast-RADIUS Attack

Attack summary: MD5 collision attack on RADIUS authentication by MITM adversary.

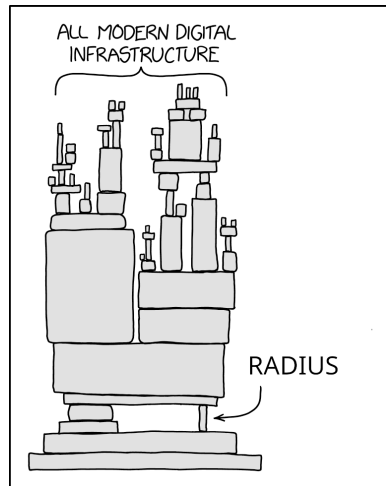


<https://blastradius.fail>

RADIUS/UDP Considered Harmful

Sharon Goldberg, Miro Haller, Nadia Heninger, Mike Milano, Dan Shumow, Marc Stevens, and Adam Suhl.

USENIX Security, August 2024.

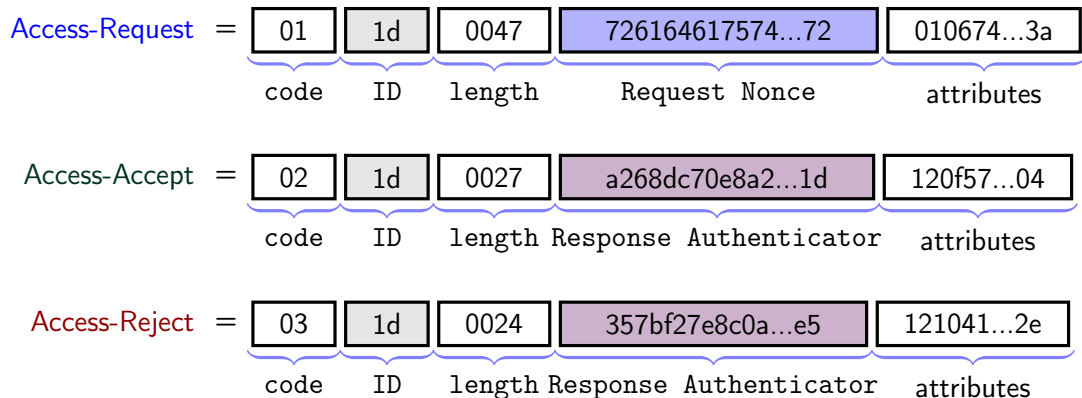


XKCD modified from [5]



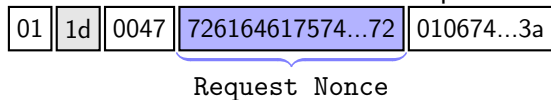
BONUS MATERIAL

End-to-End Example Attack (1/4)



End-to-End Example Attack (2/4)

1. Attacker triggers Access-Request.
2. MITM attacker observes Access-Request.



PoC example packets

`blastradius.fail/example.py`

3. MITM attacker predicts the following prefixes

Accept Prefix =

02	1d	01c0	726164617574...72	21	ec
----	----	------	-------------------	----	----

Reject Prefix =

03	1d	01c0	726164617574...72	21	ec
----	----	------	-------------------	----	----

PS (1/2)

to compute the MD5 chosen-prefix collision gibberish.

Accept Gibberish =

3d...86	21	c0	f5...9e
---------	----	----	---------

 (428 bytes)

Reject Gibberish =

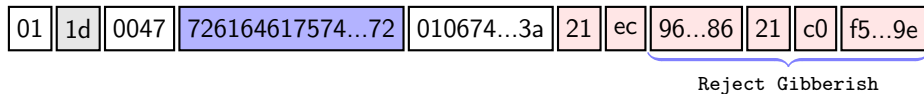
96...86	21	c0	f5...9e
---------	----	----	---------

 (428 bytes)

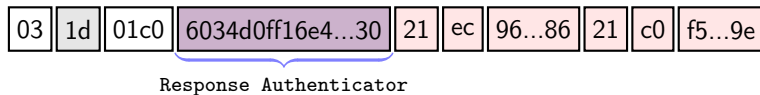
PS (2/2) Proxy State (PS)

End-to-End Example Attack (3/4)

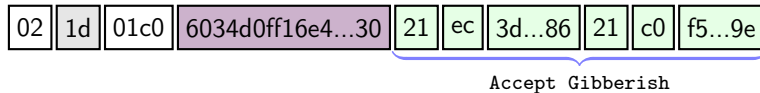
4. MITM sends Access-Request with appended Reject Gibberish to server.



5. MITM intercepts Access-Reject, learning the Response Authenticator.



6. MITM puts Response Authenticator in Access-Accept packet with appended Accept Gibberish.



End-to-End Example Attack (4/4)

7. Access-Accept and Access-Reject produce the same Response Authenticator, and hence pass the RADIUS client authentication check.

Response Authenticator

6034d0ff16e4...30

$$= \text{MD5}\left(\underbrace{\begin{array}{|c|c|c|c|} \hline 02 & 1d & 01c0 & 726164617574\dots72 \\ \hline \end{array}}_{\text{Accept Prefix}} \underbrace{\begin{array}{|c|c|c|c|c|c|} \hline 21 & ec & 3d\dots86 & 21 & c0 & f5\dots9e \\ \hline \end{array}}_{\text{Accept Gibberish}} \underbrace{\begin{array}{|c|} \hline \text{Shared Secret} \\ \hline \end{array}}_{\text{(unknown)}} \right)$$

$$= \text{MD5}\left(\underbrace{\begin{array}{|c|c|c|c|} \hline 03 & 1d & 01c0 & 726164617574\dots72 \\ \hline \end{array}}_{\text{Reject Prefix}} \underbrace{\begin{array}{|c|c|c|c|c|c|} \hline 21 & ec & 96\dots86 & 21 & c0 & f5\dots9e \\ \hline \end{array}}_{\text{Reject Gibberish}} \underbrace{\begin{array}{|c|} \hline \text{Shared Secret} \\ \hline \end{array}}_{\text{(unknown)}} \right)$$

Attack Extensions

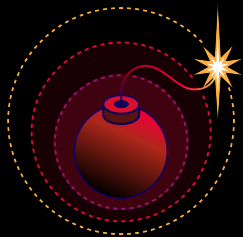
- Adversary can add arbitrary attributes in prefix for Access-Accept.

AcceptPrefix =

02	1d	01c0	726164617574...72	1a0b000007db1d04	21	ec
----	----	------	-------------------	------------------	----	----

Attribute:
Exec-Privilege 04

- Proxy-State attributes are *not* the only way to inject the RejectGibberish.
 - Any reflected user input could work, e.g. User-Name or Vendor-Specific attributes.
 - In Access-Request:
User-Name: OPZjN-_ayr83S-nc6q...Mt85
 - In Access-Reject:
Reply-Message: Login for OPZjN-_ayr83S-nc6q...Mt85 failed!
 - The client does not need to support or parse these attributes.



REFERENCES

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- [2] Alan DeKok. *Personal Communication*.
- [3] Alan DeKok. *RADIUS and MD5 Collision Attacks*. https://networkradius.com/assets/pdf/radius_and_md5_collisions.pdf. 2024.
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- [6] The FreeRADIUS Server Project and Contributors. *FreeRADIUS*. <https://freeradius.org/>.

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- [8] *SHAttered*. <https://shattered.io/>. Accessed March 20, 2025.
- [9] Marc Stevens, Arjen K. Lenstra, and Benne de Weger. “Chosen-Prefix Collisions for MD5 and Colliding X.509 Certificates for Different Identities”. In: *EUROCRYPT*. Vol. 4515. Lecture Notes in Computer Science. Springer, 2007, pp. 1–22.
- [10] Marc Stevens et al. “Short Chosen-Prefix Collisions for MD5 and the Creation of a Rogue CA Certificate”. In: *CRYPTO*. Vol. 5677. Lecture Notes in Computer Science. Springer, 2009, pp. 55–69.

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- [11] Henrik Thejl, Nagaraja K S, and Karl-Georg Aspacher. “A method for user management and a power plant control system thereof for a power plant system”. Pat. 2765466. Siemens Gamesa Renewable Energy A/S. Jan. 24, 2014. URL: <https://data.epo.org/publication-server/rest/v1.0/publication-dates/20190904/patents/EP2765466NWB1/document.pdf>.
- [12] Xiaoyun Wang and Hongbo Yu. “How to Break MD5 and Other Hash Functions”. In: *EUROCRYPT*. Vol. 3494. Lecture Notes in Computer Science. Springer, 2005, pp. 19–35.