The Blast-RADIUS Attack

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Attack Summary

A man-in-the-middle network attacker can forge arbitrary RADIUS responses for non-EAP authentication modes:

- Turning an Access-Reject into an Access-Accept
- Adding arbitrary network access attributes to Access-Accept.

This is a **protocol vulnerability**: RADIUS hard codes weak authentication based on broken MD5 hash function.

Our paper will appear at USENIX Security 2024 and is available on https://www.blastradius.fail.

Obligatory XKCD



• RADIUS is the de facto standard lightweight protocol for authentication, authorization, and accounting (AAA) for networked devices.

• RADIUS is *everywhere*: ISPs (DSL/FTTH), 802.1X, WiFi, mobile roaming, IoT, router admin access...

RADIUS Protocol Reminder



- RADIUS requests and responses are often sent over UDP.
- RFC 6614 TLS Encryption for RADIUS (2012) never left experimental status.
- (D)TLS Encryption for RADIUS: draft-ietf-radext-radiusdtls-bis

Apologies to XKCD

RADIUS Packet Formats



- Access-Request packet contents unauthenticated.
- UDP source IP address is used to identify/validate client.
- Client and server share fixed shared secret.
- Passwords obfuscated using MD5+shared secret.

RADIUS Packet Formats

Access-Accept



• Response Authenticator: Ad hoc authentication with MD5 hash.

Response Authenticator

Goal: Prevent forgery of packets, e.g., by man-in-the-middle attacker.

The Response Authenticator from packet



Blast-RADIUS: Turning Access-Reject Into Access-Accept

• Our attack allows a MITM attacker to produce a valid Response Authenticator *without* knowing the Shared Secret.

• It creates an MD5 collision such that Access-Accept and Access-Reject produce the same Response Authenticator (very simplifed):

MD5(Access-Accept) = MD5(Access-Reject).

MD5 Collision Attack History

1993 Known weaknesses in MD5.

2004 First full MD5 collision. Produced unstructured strings G_1 , G_2 with

 $\mathsf{MD5}(G_1) = \mathsf{MD5}(G_2).$

 \implies Unstructured G_1, G_2 hard to exploit in realistic contexts.

2004 Identical-prefix collision. Given prefix P, produce G_1 , G_2 such that

 $MD5(P||G_1) = MD5(P||G_2).$

 \implies Takes seconds on a laptop. Used to create colliding PDFs etc.

MD5 Collision Attack History

2007 Chosen-prefix collision. Given prefixes P_1 and P_2 , produce G_1 and G_2 such that

 $MD5(P_1||G_1) = MD5(P_2||G_2).$

Because of MD5 structure, can append any fixed suffix *S* and still collide:

 $MD5(P_1||G_1||S) = MD5(P_2||G_2||S)$

 \implies This is what we want to exploit in a real protocol.

A RADIUS Response Authenticator MD5 Collision



- 1. Given a RADIUS request and two possible responses, an attacker can add an attribute that causes them to have colliding Response Authenticators.
- 2. E.g., a forged Access-Accept and expected observed Access-Reject.
- 3. Attacker can copy valid Response Authenticator from observed response to desired forged response.

A RADIUS Response Authenticator MD5 Collision

Hence, the adversary's goal is to compute the following chosen prefix collision:

MD5 (AcceptPrefix	AcceptGibberish	Shared Secret
	=	
MD5 (RejectPrefix	RejectGibberish	Shared Secret

Problem 1: Collision prefixes AcceptPrefix and RejectPrefix depend on Request Authenticator and ID from Access Request.

Problem 2: Attacker needs to hide collision gibberish AcceptGibberish in an attribute that is echoed back from the server.

Resolving Problem 1: Online Collision Computation

Problem 1: Collision prefixes AcceptPrefix and RejectPrefix depend on Request Authenticator and ID from Access Request.

Solution: Attacker

- 1. intercepts request from victim client,
- 2. learns Request Authenticator,
- 3. predicts prefixes,
- 4. and computes MD5 collision before client timeout.

We optimized collision to get time from multiple hours to under 5 minutes for proof-of-concept attack.

We could have decreased it further by implementing on GPU/FPGA but did not think this was a good use of grad student time.

Resolving Problem 2: Reflection Via Proxy-State Attribute

Problem 2: Attacker needs to hide collision gibberish in an attribute that is echoed back from the server.

Solution: The Proxy-State attribute.

This Attribute is available to be sent by a proxy server to another server when forwarding an Access-Request and MUST be returned unmodified in the Access-Accept, Access-Reject or Access-Challenge.

(RFC 2058)

 \implies Attacker intercepts Access-Request and injects malicious Proxy-State attribute to force collision.

Blast-RADIUS Attack Overview



- 1. Adversary logs into victim NAS with invalid password.
- 2. Victim NAS makes RADIUS Access-Request to RADIUS server.
- 3. MITM intercepts request and computes MD5 collision.
- 4. MITM injects collision gibberish into Proxy-State attribute in request.
- 5. Victim RADIUS server rejects request.
- 6. MITM copies Response Authenticator from reject into forged Access-Accept.
- 7. Victim NAS RADIUS client receives forged accept and permits login.

Blast-RADIUS Attack Example (1/3)

- 1. Attacker triggers Access-Request.
- 2. MITM attacker observes Access-Request.

01 1d 0047 726164617574...72 010674...3a

Request Authenticator

3. MITM attacker predicts the following prefixes

to compute the MD5 chosen-prefix collision gibberish.

AcceptGibberish =
$$21 \text{ ec} 3d...86 21 \text{ co} \text{ f5...9e}$$
 (428 bytes)
RejectGibberish = $21 \text{ ec} 96...86 21 \text{ co} \text{ f5...9e}$ (428 bytes)
Proxy State Proxy State

PoC example packets

blastradius.fail/example.py

Blast-RADIUS Attack Example (2/3)

4. MITM sends Access-Request with appended ${\tt RejectGibberish}$ to server.



5. MITM intercepts Access-Reject, learning the Response Authenticator.

Response Authenticator

6. MITM puts Response Authenticator in Access-Accept packet with appended AcceptGibberish.

AcceptGibberish

Blast-RADIUS Attack Example (3/3)

7. Access-Accept and Access-Reject produce the same Response Authenticator, and, hence, pass the RADIUS client authentication check.

Response Authenticator





Attack Extensions

• Adversary can add arbitrary attributes in prefix for Access-Accept.

AcceptPrefix = 02 1d 01c0 726164617574...72 1a0b000007db1d04

Attribute:

Exec-Privilege 04

- Proxy-State attributes are *not* the only way to inject the RejectGibberish.
 - Any reflected user input could work, e.g. the User-Name or Vendor-Specific attributes.
 - In Access-Request:
 - User-Name: 0PZjN-_ayr83S-nc6q...Mt85
 - In Access-Reject: Reply-Message: Login for OPZjN-_ayr83S-nc6q...Mt85 failed!
 - The client does not need to support or parse these attributes.

Impact and Mitigation

Impact:

- PAP, CHAP, MS-CHAP are vulnerable
- Modes *requiring* a Message-Authenticator attribute are not vulnerable.
 - E.g., EAP.
 - HMAC-MD5 is not vulnerable to MD5 collision attack.

Threat model: Requires MITM network access

- RADIUS/UDP traffic over open internet is exposed/vulnerable.
- RADIUS/UDP traffic over VLAN or IPSEC requires attacker to have network access to exploit; useful for lateral movement within org.

Short-term mitigation: All requests and responses should include and verify Message-Authenticator attribute: draft-ietf-radext-deprecating-radius-02.

Long-term mitigation: All RADIUS traffic should be encapsulated in (D)TLS tunnel: draft-ietf-radext-radiusdtls-bis-02.

Discussion

MD5 has been known to have weaknesses for 30 years.

MD5 has been cryptographically broken for 20 years.

The absence of an explicit attack against MD5 in a given protocol likely indicates cryptographers are unaware it is still in use, rather than security.

Shout out to CERT and Alan for coordinating disclosure process.

Blast-RADIUS Takeaways

Attack summary: Man-in-the-middle attack allowing arbitrary forged RADIUS responses for non-EAP authentication methods.

Long-term mitigation: Standardize and move to RADIUS/TLS.

RADIUS/UDP Considered Harmful

Sharon Goldberg, Miro Haller, Nadia Heninger, Mike Milano, Dan Shumow, Marc Stevens, and Adam Suhl. To appear at USENIX Security, August 2024.

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